

Clarification / correction:

C_n & h_n are defined for $n \geq 3$
 k_n defined for $n \geq 1$

$k_1: 0$ $k_2: 0 \rightarrow 0$

Recursion

recursive function defn has 2 parts:

- ① base case(s)
- ② recursive formula

Fibonacci numbers:

$$\left. \begin{array}{l} F_0 = 0 \\ F_1 = 1 \end{array} \right\} \text{base cases}$$

$$F_i = F_{i-1} + F_{i-2} \quad \forall i \geq 2$$

$$F_5? = F_4 + F_3$$

$$(F_3 + F_2) + (F_2 + F_1)$$

$$(F_2 + F_1) + (F_1 + F_0) + (F_1 + F_0) + F_1$$

$$F_1 + F_0 + F_1 + F_1 + F_0 + F_1 + F_0 + F_1$$

$$1 + 0 + 1 + 1 + 0 + 1 + 0 + 1 = 5$$

$$\begin{array}{ccccccc} 0, & 1, & 1, & 2, & 3, & 5 & \text{"normal way"} \\ \downarrow & & & & & \downarrow & \\ 0^{\text{th}} & & & & & 5^{\text{th}} & \end{array}$$

unrolling, since my input = 5
I completed unrolling up to
base case

unrolling: technique to find closed form of a recursive function
not a proof technique.
induction to prove it.

example: $T: \mathbb{N} \rightarrow \mathbb{Z}$

$$T(1) = 1$$

$$T(n) = 2T(n-1) + 3$$

level

$$k=1 \quad T(n) = 2T(n-1) + 3$$

$$k=2 \quad T(n) = 2(2T(n-2) + 3) + 3$$

$$k=3 \quad T(n) = 2(2(2T(n-3) + 3) + 3) + 3$$

$$k=3 \quad T(n) = 2^3 T(n-3) + 2^2 \cdot 3 + 2^1 \cdot 3 + 2^0 \cdot 3$$

$$k? \quad T(n) = 2^k T(n-k) + \sum_{i=0}^{k-1} 2^i \cdot 3 \quad (1)$$

$$= T(1)$$

$$n-k=1 \rightarrow k=n-1 \quad (2)$$

value of k at base case of 1

$$2^{n-1} T(n-(n-1)) + \sum_{i=0}^{n-2} 2^i \cdot 3$$

$$2^{n-1} \cdot \underline{T(1)} + 3 \sum_{i=0}^{n-2} 2^i$$

$$2^{n-1} + \underline{3 \sum_{i=0}^{n-2} 2^i}$$

$$2^{n-1} + 3(2^{n-1} - 1)$$

$$2^{n-1} + 3 \cdot 2^{n-1} - 3$$

$$4 \cdot 2^{n-1} - 3$$

$$2^2 \cdot 2^{n-1} - 3$$

$$\underline{2^{n+1} - 3}$$

closed form

example: $S(1) = C$
 $S(n) = 2S(n/2) + n \quad \forall n \geq 2$

$$S(6) = 2S(3) + \underline{6}$$
$$S(n/2) = 2S(n/4) + \underline{n/2} \quad \star \text{ don't forget}$$

level

$$k=1 \quad S(n) = 2S(n/2) + n$$

$$k=2 \quad S(n) = 2(2S(n/4) + n/2) + n$$

$$k=3 \quad S(n) = \underline{2}(\underline{2}(2S(n/8) + \underline{n/4}) + \underline{n/2}) + n$$

$$k \quad S(n) = \underline{2^k S(n/2^k)} + kn$$

$$k? \quad S(1)$$

$$n/2^k = 1$$

$$n = 2^k$$

$$k = \log_2 n$$

$$S(n) = 2^{\log_2 n} \cdot S(n/2^{\log_2 n}) + (\log_2 n)(n)$$

$$S(n) = n \cdot S(1) + n \log_2 n$$

$$S(n) = n \cdot C + n \log_2 n \quad \text{closed form}$$

\star Do 12.2 bc instead of 12.2 ab